



# The Möbius transform on symmetric ordered structures and its application to capacities on finite sets

Michel Grabisch

## ► To cite this version:

Michel Grabisch. The Möbius transform on symmetric ordered structures and its application to capacities on finite sets. *Discrete Mathematics*, 2004, 287 (1-3), pp.17-34. 10.1016/j.disc.2004.05.013 . hal-00188158

**HAL Id: hal-00188158**

**<https://hal.science/hal-00188158>**

Submitted on 15 Nov 2007

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# The Möbius transform on symmetric ordered structures and its application to capacities on finite sets

Michel GRABISCH<sup>1</sup>

*Université Paris I Panthéon-Sorbonne*

---

## Abstract

Considering a linearly ordered set, we introduce its symmetric version, and endow it with two operations extending supremum and infimum, so as to obtain an algebraic structure close to a commutative ring. We show that imposing symmetry necessarily entails non associativity, hence computing rules are defined in order to deal with non associativity. We study in details computing rules, which we endow with a partial order. This permits to find solutions to the inversion formula underlying the Möbius transform. Then we apply these results to the case of capacities, a notion from decision theory which corresponds, in the language of ordered sets, to order preserving mappings, preserving also top and bottom. In this case, the solution of the inversion formula is called the Möbius transform of the capacity. Properties and examples of Möbius transform of sup-preserving and inf-preserving capacities are given.

---

## 1 Introduction

We consider a linearly ordered set  $(L^+, \leq)$ , with bottom and top denoted by  $\mathbb{0}, \mathbb{1}$  respectively, and we define  $L := L^+ \cup L^-$ , where  $L^-$  is a reversed copy of  $L^+$ , i.e. for any  $a, b \in L^+$ , we have  $a \leq b$  iff  $-b \leq -a$ , where  $-a, -b$  are the copies of  $a, b$  in  $L^-$ . The set of signed integers, the set of real numbers have this structure, with a central zero, and they possess rich algebraic structures (groups, rings, etc.) when endowed with usual arithmetical operations.

Our aim is to build similar structures, but using only the order relation  $\leq$  on  $L$ , so that the resulting structure should be as close as possible to e.g. the

---

*Email address:* `michel.grabisch@lip6.fr` (Michel GRABISCH).

<sup>1</sup> Mailing address: LIP6 — 8, rue du Capitaine Scott, 75015 Paris, France

ring of real numbers with  $+$ ,  $\times$ . This should permit an easy manipulation of functions or functionals (such as set functions, capacities, integrals) taking values in  $L$ .

Our work is essentially motivated by decision making. Since this is fundamental for our approach, we briefly introduce necessary notions. The aim of decision making is to rank or assign overall scores to *alternatives*, i.e. functions  $f : S \longrightarrow L$ , where  $S$  is a set of features (criteria, points of view, states of nature, etc.), and  $f(s) \in L$  for any  $s \in S$  is the score of  $f$  for feature  $s$ , expressed on some scale (usually a real interval). Then overall scoring of  $f$  can be viewed as a functional  $V : L^S \longrightarrow L$  satisfying certain properties. A very general way to define  $V$  is to take some integral. The Choquet integral [4], generalizing the Lebesgue integral, has proven to be a suitable and very general functional for decision making [24], defined for non negative functions. The Choquet integral is defined with respect to a *capacity*  $v : 2^S \longrightarrow [0, 1]$ , a monotone set function extending classical measures used in the Lebesgue integral. For any capacity  $v$ , its *Möbius transform*  $m^v : 2^S \longrightarrow \mathbb{R}$  is a key concept in decision analysis (see e.g. [3]). It is defined by

$$m^v(A) := \sum_{B \subseteq A} (-1)^{|A \setminus B|} v(B), \quad \forall A \subseteq S. \quad (1)$$

In a more general way, the Möbius transform provides an inversion formula useful in combinatorics [23].

Let us denote by  $\mathcal{C}_v(f)$  the Choquet integral of  $f$  with respect to capacity  $v$ . When  $L$  happens to be a real interval containing negative numbers, then a suitable extension of the Choquet integral has to be defined for real-valued functions. It is called the *symmetric Choquet integral*, and is defined as:

$$\check{\mathcal{C}}_v(f) := \mathcal{C}_v(f^+) - \mathcal{C}_v(f^-) \quad (2)$$

with  $f^+ = f \vee 0$ , and  $f^- = (-f)^+$ . This is the basis for *Cumulative Prospect Theory* [27], an important theory in economics for representing human behaviour in decision making when faced with gains (positive values,  $L^+$ ) and losses (negative values,  $L^-$ ).

If  $L$  is only an *ordinal scale*, i.e. a (usually finite) scale with only a total order on it, then the Choquet integral is no more applicable, since usual arithmetical operations are not defined on  $L$ . It is known that the counterpart of it is the Sugeno integral [26], denoted  $\mathcal{S}_v(f)$ , which is defined solely with supremum ( $\vee$ ) and infimum ( $\wedge$ ), and like the Choquet integral, with respect to a capacity  $v$ , which has to be valued on  $L$ . However, in the ordinal case, there is no symmetric Sugeno integral, since first of all there is no notion of “negative numbers” for ordinal scales. Similarly, there is no Möbius transform. Our aim is precisely to define negative ordinal quantities so as to obtain a sufficiently

rich algebraic structure on  $L$  to allow computation similar to (1), (2), and thus to be able to develop an ordinal counterpart of Cumulative Prospect Theory.

Generally speaking, we may think of several ways to tackle this problem. We denote  $\oplus, \otimes$  the new operations on  $L$ .

An immediate solution would have been to use the Boolean ring associated to  $(L, \leq)$ . But this solution works only if  $(L, \leq)$  is a Boolean lattice, and our application field requires that  $L$  be only a linear lattice.

A second solution is to define  $\oplus, \otimes$  as binary operators  $L^2 \longrightarrow L$ , and impose (possibly among other conditions) that

- (C1)  $\oplus, \otimes$  coincide with  $\vee, \wedge$  respectively on  $L^+$
- (C2)  $-a$  is the symmetric of  $a$ , i.e.  $a \oplus (-a) = \mathbb{O}$ .
- (C3)  $-(a \oplus b) = (-a) \oplus (-b)$ ,  $-(a \otimes b) = (-a) \otimes b$ .

These conditions are motivated by our aim to develop an ordinal Cumulative Prospect Theory:

- (1) (C1) permits us to perform an extension of all that already exists in  $L^+$ , e.g. the Sugeno integral.
- (2) Thanks to (C2) and (C3), computations could be conducted as with real numbers, with  $\oplus, \otimes$  playing the role of  $+, \times$  respectively. In particular, it would permit to define a counterpart of the Möbius transform (1), and to define a *symmetric Sugeno integral*, in the spirit of (2):

$$\check{\mathcal{S}}_v(f) := \mathcal{S}_v(f^+) \oplus (-\mathcal{S}_v(f^-)) \quad (3)$$

Condition (C2) then implies that the integral of  $f$  (overall scoring) is  $\mathbb{O}$  whenever  $\mathcal{S}_v(f^+) = \mathcal{S}_v(f^-)$ , a desirable property since it means that the overall scoring should be null when gains equal losses.

The problem with this solution is that due to (C1) and (C2), inevitably  $\oplus$  would be non associative in general. Take  $\mathbb{O} < a < b$  and consider the expression  $(-b) \oplus b \oplus a$ . Depending on the place of parentheses, the result differs since  $((-b) \oplus b) \oplus a = \mathbb{O} \oplus a = a$ , but  $(-b) \oplus (b \oplus a) = (-b) \oplus b = \mathbb{O}$ . In other words, if we want to keep associativity and (C1), then necessarily, (C2) does not hold: Prop. 5 will show that in this case,  $|a \oplus (-a)| \geq |a|$ . Clearly, this result does not match intuition in our decision making perspective, and hence we have no other way than to accept non associativity. Remark however that as far as Eq. (3) is concerned, we need no associativity property.

In order to escape the incompatibility between symmetry and associativity, a third solution would be to define  $\oplus, \otimes$  as binary operators on  $(L^+ \times L^-)^2 \longrightarrow L^+ \times L^-$ , that is, each element  $a$  in  $L$  is viewed as a pair  $(a^+, a^-) \in L^+ \times L^-$ ,

where  $a^- = \mathbb{O}$  if  $a \geq \mathbb{O}$ , and  $a^+ = \mathbb{O}$  otherwise. Then for any  $(a^+, a^-), (b^+, b^-)$  in  $L^+ \times L^-$ , one could define in an obvious way:

$$\begin{aligned} -(a^+, a^-) &:= (-a^-, -a^+) \\ (a^+, a^-) \oplus (b^+, b^-) &:= (a^+ \vee b^+, a^- \wedge b^-) \\ (a^+, a^-) \otimes (b^+, b^-) &:= (a^+ \wedge b^+, a^- \vee b^-). \end{aligned}$$

Thus  $\oplus, \otimes$  are associative since  $\vee, \wedge$  are on  $L^+, L^-$ , they coincide with  $\vee, \wedge$  on  $L^+$  (condition **(C1)**), and condition **(C3)** is fulfilled since  $-((a^+, a^-) \oplus (b^+, b^-)) = (-a^+, a^-) \oplus (-b^+, b^-)$ . However, elements have no opposite, since  $(a^+, a^-) \oplus (-a^-, -a^+) \neq (\mathbb{O}, \mathbb{O})$ . Also, there is no total order on  $L^+ \times L^-$ .

Considering our motivation, only the second solution is acceptable, since symmetry is mandatory in our framework, and the third solution would lead to a partial order on alternatives, a situation which is not desirable in decision making.

In this paper, our aim is to completely develop the second solution, and to apply it in particular to the definition of an ordinal Möbius transform, the definition of the symmetric Sugeno integral being already solved as indicated above. First attempts at defining the Möbius transform of capacities in an ordinal context have been done by Marichal [18,16], Mesiar [19], and the author [9]. However, these preliminary works have been done without explicit connection to combinatorics, and were restricted to capacities.

The paper is organized as follows. Next section introduces necessary concepts for the sequel, while Section 3 gives the construction of the symmetric ordered structure. Since  $\oplus$  is necessarily non associative, Section 4 introduces rules of computation, which give meaning to expression such as  $\oplus_{i \in I} a_i$ . Section 5 is devoted to the study of the Möbius inversion formula, when defined on symmetric ordered structures. Lastly, Section 6 focusses on capacities, while Section 7 concludes the paper by indicating possible applications of the results.

## 2 Preliminaries

We give necessary definitions and introduce basic concepts for our construction.

Let  $N$  be a finite set, and  $(L^+, \leq)$  a totally ordered set, with  $\mathbb{O}, \mathbb{1}$  being top and bottom elements. A  $(L^+)$ -valued *capacity* is an order preserving (or *isotone*) mapping  $v : (2^N, \subseteq) \longrightarrow (L^+, \leq)$ , with  $v(\emptyset) = \mathbb{O}$ ,  $v(N) = \mathbb{1}$ .

We say that a complete lattice  $(\mathcal{L}, \leq)$  is a *conjugation* lattice if it is endowed with a bijective and order-reversing mapping from  $\mathcal{L}$  to  $\mathcal{L}$ , called a *conjugation*, which maps  $a$  to  $\bar{a}$ , so that  $\bar{\bar{a}} = a$ , and  $a \leq b$  iff  $\bar{a} \geq \bar{b}$ . Then  $\overline{a \vee b} = \bar{a} \wedge \bar{b}$ , and  $\overline{a \wedge b} = \bar{a} \vee \bar{b}$ . In the Boolean lattice  $2^N$ , set complement is a conjugation.

If  $(L^+, \leq)$  has a conjugation, then the *conjugate capacity*  $\bar{v}$  is defined by  $\bar{v}(A) := \overline{v(\bar{A})}$ ,  $A \subseteq N$ .

Let us consider a poset  $(P^+, \leq)$ , with bottom and top elements denoted by  $\mathbb{0}$  and  $\mathbb{1}$ . We introduce  $P^- := \{-a | a \in P^+\}$ , with the reversed order, i.e.  $-a \leq -b$  iff  $a \geq b$  in  $P^+$ . The bottom and top of  $P^-$  are respectively  $-\mathbb{1}$  and  $-\mathbb{0}$ .

The disjoint union of  $P^+$  and  $P^-$ , with identification of  $-\mathbb{0}$  with  $\mathbb{0}$ , is called a *reflection poset* or *symmetric poset*, and is denoted  $(P, \leq)$  [5]. It is a poset with bottom  $-\mathbb{1}$  and top  $\mathbb{1}$ .

We introduce some mappings on  $(P, \leq)$ . The *reflection* maps  $a \in P$  to  $-a$ , and  $-(-a) := a$  for any  $a \in P$ . If  $P$  is a lattice we have:

$$(-a) \vee (-b) = -(a \wedge b), \quad (-a) \wedge (-b) = -(a \vee b).$$

The *absolute value* of  $a \in P$  is denoted  $|a|$ , and  $|a| := a$  if  $a \in P^+$ , and  $|a| = -a$  otherwise. The *sign function* is defined by:

$$\text{sign} : P \rightarrow P, \quad \text{sign } x = \begin{cases} -\mathbb{1} & \text{for } x < \mathbb{0} \\ \mathbb{0} & \text{for } x = \mathbb{0} \\ \mathbb{1} & \text{for } x > \mathbb{0} \end{cases}.$$

Lastly, we introduce the notion of derivative.

**Definition 1** Let  $(X, \leq)$  be a poset and  $(\mathcal{L}, \leq)$  be a complete lattice. For any isotone function  $g$  from  $X$  to  $\mathcal{L}$ , let  $\tilde{g}$  be defined by:

$$\tilde{g}(x) := \bigvee_{y < x} g(y). \tag{4}$$

Then the derivative  $g'$  of  $g$  is defined by:

$$g'(x) := \begin{cases} \mathbb{0}, & \text{if } g(x) = \tilde{g}(x), \\ g(x), & \text{otherwise.} \end{cases}$$

In a partially ordered set (poset for short)  $(X, \leq)$ , we say that  $x$  *covers*  $y$ , denoted by  $x \succ y$ , if  $x > y$ , and there is no  $u \in X$  such that  $u \neq x, y$  and  $x > u > y$ .

### 3 Symmetric ordered structures

Let  $(L^+, \leq)$  be a totally ordered set (linear lattice) with top and bottom  $\mathbb{1}, \mathbb{0}$ , and consider the corresponding symmetric linear lattice  $(L, \leq)$ . As stated in the introduction, our aim is to endow  $L$  with two operations denoted  $\oplus, \otimes$  extending usual  $\vee, \wedge$  on  $L^+$ , so that the resulting structure is close to a ring. More precisely, we require the following:

- (C1)  $\oplus, \otimes$  coincide with  $\vee, \wedge$  respectively on  $L^+$
- (C2)  $-a$  is the symmetric of  $a$ , i.e.  $a \otimes (-a) = \mathbb{0}$ ,  $\forall a \in L$ .
- (C3)  $-(a \otimes b) = (-a) \otimes (-b)$ ,  $-(a \oplus b) = (-a) \oplus b$ ,  $\forall a, b \in L$ .

Let us build  $\otimes$  first. Due to (C1),  $\otimes$  is defined on  $L^+$  and coincides with  $\vee$ , which implies that  $a \otimes \mathbb{0} = a$  for any  $a > \mathbb{0}$ . Using (C3), we deduce  $-(a \otimes \mathbb{0}) = (-a) \otimes \mathbb{0} = -a$ , showing that  $\mathbb{0}$  is the neutral element. Again using (C3) with  $a, b \geq \mathbb{0}$ , we get

$$-(a \otimes b) = -(a \vee b) = (-a) \wedge (-b) = (-a) \otimes (-b)$$

showing that  $\otimes$  coincides with  $\wedge$  on  $L^-$ . It remains to define  $\otimes$  for arguments of opposite sign. We propose the following (this will be justified in Prop. 5), assuming  $a > \mathbb{0}$  and  $b < \mathbb{0}$

$$a \otimes b := \begin{cases} a, & \text{if } a > -b \\ \mathbb{0}, & \text{if } a = -b \\ b, & \text{otherwise.} \end{cases} \quad (5)$$

Note that the second case is just (C2). Using (C3), we can derive the formula for the opposite case  $a < \mathbb{0}$  and  $b > \mathbb{0}$ . In summary,  $\otimes$  is given by Fig. 1.

A compact formulation of  $\otimes$  is:

$$a \otimes b := \begin{cases} -(|a| \vee |b|) & \text{if } b \neq -a \text{ and } |a| \vee |b| = -a \text{ or } = -b \\ \mathbb{0} & \text{if } b = -a \\ |a| \vee |b| & \text{else.} \end{cases} \quad (6)$$

Except for the case  $b = -a$ ,  $a \otimes b$  equals the larger one in absolute value of the two elements  $a$  and  $b$ .

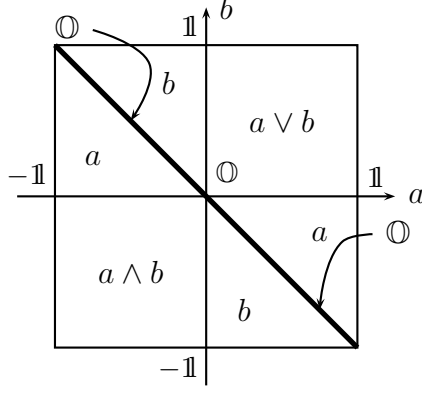


Fig. 1. Definition of the symmetric maximum  $a \odot b$

**Remark 2** The following interpretation can be given for  $\odot$ : on scale  $L$ , distinct levels are far away from one another, so that invoking negligibility aspects, only the maximum level remains when combining two positive values. When a positive (gain) and a negative value (loss) are combined, if the gain dominates the loss, the latter counts for nothing.

**Remark 3** Equation (5) is a symmetrized version of a difference operator introduced by Weber [28]:

$$a \sharp b := \inf\{c | b \vee c \geq a\} = \begin{cases} a, & \text{if } a > b \\ \odot, & \text{otherwise} \end{cases} \quad (7)$$

for any  $a, b \in L^+$ . Note that  $a \sharp b$  is the dual of the pseudo-complement of  $b$  relative to  $a$ , defined by  $b * a := \sup\{c | b \wedge c \leq a\}$  (see e.g. [13]). It is also the dual of the residual of  $a$  by  $b$  (see e.g. [2]).

It remains to define the *symmetric minimum* operator. Since we impose the symmetry condition (**C3**), we are naturally lead to Fig 2.

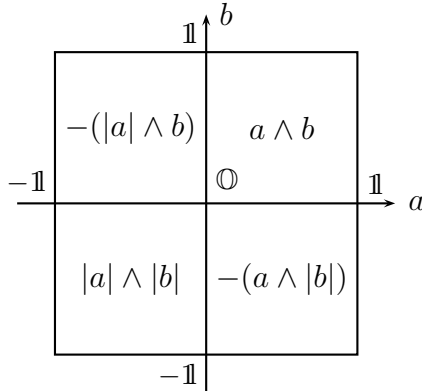


Fig. 2. Definition of the symmetric minimum  $a \otimes b$



A more compact expression is:

$$a \oslash b := \begin{cases} -(|a| \wedge |b|) & \text{if } \text{sign } a \neq \text{sign } b \\ |a| \wedge |b| & \text{else.} \end{cases} \quad (8)$$

The absolute value of  $a \oslash b$  equals  $|a| \wedge |b|$  and  $a \oslash b < \mathbb{O}$  iff the two elements  $a$  and  $b$  have opposite signs.

Another equivalent formulation of these two operations, applicable if  $L$  is a symmetric real interval, is due to Marichal [17], and clearly shows the relationship with the ring of real numbers.

$$a \oslash b = \text{sign}(a + b)(|a| \vee |b|) \quad (9)$$

$$a \oslash b = \text{sign}(a \cdot b)(|a| \wedge |b|). \quad (10)$$

The following proposition summarizes the properties of the structure obtained.

**Proposition 4** *The structure  $(L, \oslash, \mathbb{O})$  has the following properties.*

- (i)  $\oslash$  is commutative.
- (ii)  $\mathbb{O}$  is the unique neutral element of  $\oslash$ , and the unique absorbing element of  $\oslash$ .
- (iii)  $a \oslash -a = \mathbb{O}$ , for all  $a \in L$ .
- (iv)  $-(a \oslash b) = (-a) \oslash (-b)$ .
- (v)  $\oslash$  is associative for any expression involving  $a_1, \dots, a_n$ ,  $a_i \in L$ , such that  $\bigvee_{i=1}^n a_i \neq -\bigwedge_{i=1}^n a_i$ .
- (vi)  $\mathbb{O}$  is commutative.
- (vii)  $\mathbb{1}$  is the unique neutral element of  $\oslash$ , and the unique absorbing element of  $\oslash$ .
- (viii)  $\oslash$  is associative on  $L$ .
- (ix)  $\oslash$  is distributive w.r.t  $\oslash$  in  $L^+$  and  $L^-$  separately.
- (x)  $\oslash$  is isotone, i.e.  $a \leq a', b \leq b'$  implies  $a \oslash b \leq a' \oslash b'$ .

**PROOF.** All results are almost clear from the construction. We just detail (v) and (ix).

(v) Let us study if the equality  $(a \oslash b) \oslash c = a \oslash (b \oslash c)$  holds supposing there is no pair of symmetric elements, as  $(a, -a)$ . This implies  $|a \oslash b| = |a| \vee |b|$  (see (6)). Hence

$$|(a \oslash b) \oslash c| = |a \oslash b| \vee |c| = |a| \vee |b| \vee |c| = |a| \vee |b \oslash c| = |a \oslash (b \oslash c)|.$$

Thus,  $(a \oslash b) \oslash c$  and  $a \oslash (b \oslash c)$  have the same absolute value. It remains to prove that they have the same sign. The sign of  $a \oslash b$  is the sign of the largest term in absolute value. Hence, the sign of  $(a \oslash b) \oslash c$  is the sign of the largest in absolute value among  $a \oslash b$  and  $c$ , so it is the sign of the largest in absolute value among  $a, b, c$ . Doing the same with  $a \oslash (b \oslash c)$ , we conclude that the two expressions have the same sign.

Suppose now  $a = -b$ . Then  $(a \oslash b) \oslash c = c$ . Clearly,  $a \oslash (b \oslash c) = c$  if and only if  $|c| > |a|$ . This coincides with the condition given in (v).

(ix) Distributivity is clearly satisfied on  $L^+$ . For any  $a, b, c \in L^+$ :

$$\begin{aligned} (a \oslash b) \oslash c &= (a \wedge b) \oslash |c| = |a \wedge b| \wedge |c| \\ (a \oslash c) \oslash (b \oslash c) &= (|a| \wedge |c|) \vee (|b| \wedge |c|) \\ &= (|a| \vee |b|) \wedge |c| \\ &= |a \wedge b| \wedge |c|. \end{aligned}$$

□

The distributivity does not hold in general: take  $a, b \geq \mathbb{O}$ ,  $a < b$ ,  $c < \mathbb{O}$ ,  $b < -c$ . Then

$$\begin{aligned} a \oslash (b \oslash c) &= a \oslash c = -a \\ (a \oslash b) \oslash (a \oslash c) &= a \oslash (-a) = \mathbb{O}. \end{aligned}$$

Using the definition of the symmetric maximum, we see that the derivative of a function  $g$  (see Definition 1) can be reformulated as:

$$g'(x) = g(x) \oslash (-\tilde{g}(x)). \quad (11)$$

The next proposition gives justifications to our choice in (5), and of the overall construction.

**Proposition 5** *We consider conditions (C1), (C2) and (C3), and denote by (C3+) condition (C3) when  $a, b$  are restricted to  $L^+$ . Then:*

(0) *Conditions (C1) and (C2) implies that associativity cannot hold.*

(1) *Under conditions (C1), (C2) and (C3), no operation is associative on a larger domain than  $\oslash$  as given by (6).*

(2) *Under (C1) and (C3+),  $\mathbb{O}$  is neutral. If we require in addition associativity, then  $|a \oslash (-a)| \geq |a|$ . Further, if we require isotonicity of  $\oslash$ , then  $|a \oslash (-a)| = |a|$ .*

**PROOF.** (0) Let us take  $\mathbb{O} < a < b$ . Then  $((-b) \otimes b) \otimes a = \mathbb{O} \otimes a = a \neq (-b) \otimes (b \otimes a) = (-b) \otimes b = \mathbb{O}$  (see introduction).

(1) The only degree of freedom is the definition of  $a \otimes b$  when  $a, b$  have different signs. We know that the only non associative case happens in expressions like  $-x \otimes (x \otimes y)$ ,  $x, y \geq \mathbb{O}$ . Since  $\otimes \equiv \vee$  on  $[\mathbb{O}, \mathbb{1}]^2$ , we get:

$$-x \otimes (x \otimes y) = -x \otimes (x \vee y) = \begin{cases} -x \otimes x = \mathbb{O}, & \text{if } x \geq y \\ -x \otimes y, & \text{if } x \leq y. \end{cases}$$

Observing that  $(-x \otimes x) \otimes y = y$ , clearly the first case can never lead to associativity. Let us examine the second case. It leads to associativity iff  $-x \otimes y = y$ . Discarding the case  $x = y$ , we see that we have in fact the definition of the symmetric maximum. Hence only it can lead to associativity in this case, and any other operation would not.

(2) Let us assume **(C1)** and **(C3+)**. If  $a > \mathbb{O}$ , then  $a \otimes \mathbb{O} = a$ , and  $-(a \otimes \mathbb{O}) = (-a) \otimes \mathbb{O} = -a$ . Now, if associativity holds, then taking  $a > \mathbb{O}$ , we have  $((-a) \otimes a) \otimes a = (-a) \otimes (a \otimes a)$ , which gives  $((-a) \otimes a) \otimes a = (-a) \otimes a$ . We know from (0) that **(C1)** and associativity imply that  $(-a) \otimes a \neq \mathbb{O}$ . If  $(-a) \otimes a > \mathbb{O}$ , then to satisfy the above equality we must have  $(-a) \otimes a \geq a$ . If it is a negative, a similar argument using **(C3+)** shows that  $(-a) \otimes a \leq -a$ . Lastly, if  $\otimes$  is isotone, we have  $a \otimes (-a) \leq a \otimes \mathbb{O} = a$ , and similarly for the negative case.  $\square$

#### 4 Non associativity and computing rules

Due to the lack of associativity of  $\otimes$ , expressions like  $\otimes_{i=1}^n a_i$  have no meaning, unless one defines a particular and systematic way of arranging terms so that associativity problems disappear.

Let us consider a sequence  $\{a_i\}_{i \in I}$  of terms  $a_i \in L$ , with  $I \subseteq \mathbb{N}$ . We say that the sequence *fulfills associativity* if either  $|I| \leq 2$  or  $\vee_{i \in I} a_i \neq \wedge_{i \in I} a_i$ . Hence, from Prop. 4 (v),  $\otimes_{i \in I} a_i$  is well-defined if and only if the sequence  $\{a_i\}_{i \in I}$  is associative. If a sequence does not fulfill associativity, it necessarily has at least 3 terms and contains a pair of maximal opposite terms  $(a, -a)$ , with  $a := \vee_{i \in I} a_i$ . Discarding all occurrences of  $a, -a$  in the sequence, we may still find (new) maximal opposite terms  $b, -b$ , which can be discarded, etc., until no more such terms remain, which means that the new sequence fulfills associativity. We call the *sequence of maximal opposite terms* the sequence of all deleted terms, whose index set is denoted  $I_-$ . Taking for example with  $L = \mathbb{Z}$  the sequence  $3, 3, 3, 2, 1, 0, -2, -3, -3$ , the sequence of maximal opposite terms is  $3, 3, 3, 2, -2, -3, -3$ .

Another way to fulfill associativity is obtained by discarding in the sequence the pair  $(a, -a)$ , with  $a := \vee_{i \in I} a_i$ , and if the new sequence  $\{a_i\}_{i \in I} \setminus \{a, -a\}$  does not fulfill associativity, then discard the pair of maximal opposite terms in this new sequence, etc., until associativity is fulfilled. We call the *restricted sequence of maximal opposite terms* the sequence of all deleted terms, and we denote its index set by  $I_0$ . In the previous example, the restricted sequence of maximal opposite terms is  $3, -3, 3, -3$ . Note that we always have  $I_0 \subseteq I_-$ , and that  $I_0$  is minimal in the sense that no proper subset of it can ensure associativity.

We denote the set of all (at most countable) sequences, including the empty one, by  $\mathfrak{S} := \bigcup_{i=1}^{\infty} L^i \cup \{\emptyset\}$ . From now on, we make the convention  $\mathbb{Q}_\emptyset a_i = \mathbb{Q}$ .

A *computation rule* is a systematic way to delete terms in a sequence  $\{a_i\}_{i \in I}$ , so that it fulfills associativity, provided the way they are deleted can be obtained as the result of a suitable arrangement of parentheses in  $\mathbb{Q}_{i \in I} a_i$ . For example, deleting 3 in the sequence  $3, 1, -3$  makes the sequence associative, but does not correspond to some arrangement of parentheses, and so is not a computation rule. Formally, we denote a computation rule by the infix notation:

$$\begin{aligned} \langle \cdot \rangle : \mathfrak{S} &\longrightarrow \mathfrak{S} \\ \{a_i\}_{i \in I} &\mapsto \langle \{a_i\}_{i \in I} \rangle := \{a_i\}_{i \in I \setminus J} \end{aligned}$$

where  $J \subseteq I$  is the index set of deleted terms. To avoid heavy notations, we denote  $\mathbb{Q}_{i \in I \setminus J} \langle \{a_i\}_{i \in I} \rangle$  by  $\langle \mathbb{Q}_{i \in I} a_i \rangle$ . The set of all computation rules on  $L$  is denoted by  $\mathfrak{R}$ .

Let us give some basic examples of computation rule.

- (1) The *weak rule*  $\langle \cdot \rangle_=-$ , where the index set of deleted terms is  $J = I_-$ . It obviously corresponds to a particular arrangement of parentheses, as shown in the following example:

$$\begin{aligned} \langle 3 \otimes 3 \otimes 3 \otimes 2 \otimes 1 \otimes 0 \otimes -2 \otimes -3 \otimes -3 \rangle_- = \\ ((3 \otimes 3 \otimes 3) \otimes (-3 \otimes -3)) \otimes (2 \otimes -2) \otimes (1 \otimes 0) = 1. \end{aligned} \quad (12)$$

- (2) The *strong rule*  $\langle \cdot \rangle_0$ , whose index set of deleted terms is  $I_0$ . It obviously corresponds to a particular arrangement of parentheses. Our example gives

$$\begin{aligned} \langle 3 \otimes 3 \otimes 3 \otimes 2 \otimes 1 \otimes 0 \otimes -2 \otimes -3 \otimes -3 \rangle_0 = \\ (3 \otimes -3) \otimes (3 \otimes -3) \otimes (3 \otimes 2 \otimes 1 \otimes 0 \otimes -2) = 3. \end{aligned} \quad (13)$$

- (3) The *splitting rule*  $\langle \cdot \rangle_+^+$ , whose index set of deleted terms is  $J = \emptyset$  if the sequence fulfills associativity, and  $J = I$  if not. Then in the latter case,

$\langle \bigotimes_{i \in I} a_i \rangle_-^+ = \mathbb{O}$ , due to our convention  $\bigotimes_{\emptyset} a_i = 0$ . The corresponding arrangement of parentheses is

$$\langle \bigotimes_{i \in I} a_i \rangle_-^+ := \left( \bigotimes_{a_i \geq \mathbb{O}} a_i \right) \bigotimes \left( \bigotimes_{a_i < \mathbb{O}} a_i \right).$$

hence the name of the rule (splitting positive and negative terms).

- (4) The *optimistic* and *pessimistic* rules  $\langle \cdot \rangle_{\text{opt}}$ ,  $\langle \cdot \rangle_{\text{pes}}$ . Let us consider a sequence of at least 3 terms in  $\mathfrak{S}$  having maximal opposite terms  $a, -a$ , with degrees of multiplicity  $k_+, k_-$  respectively. If  $k_+ = 1$  and  $k_- \leq 2$ , or  $k_- = 1$  and  $k \leq 2$ , then  $J = I$  for both rules (hence they give  $\mathbb{O}$  as result). Otherwise, the optimistic rule deletes  $k_+ - 1$  occurrences of  $a$  and all  $k_-$  occurrences of  $-a$  (hence it returns  $a$ ), while the pessimistic rule deletes all  $k_+$  occurrences of  $a$  and  $k_- - 1$  occurrences of  $-a$  (hence it returns  $-a$ ). One can verify that these rules can be expressed as a particular arrangement of parentheses. For example

$$\begin{aligned} \langle 3 \otimes 3 \otimes 3 \otimes 2 \otimes 1 \otimes 0 \otimes -2 \otimes -3 \otimes -3 \rangle_{\text{pes}} = \\ ((3 \otimes 3 \otimes 3) \otimes -3) \otimes (-3 \otimes 2 \otimes -2 \otimes 1 \otimes 0) = -3. \end{aligned} \quad (14)$$

**Remark 6** *The first three rules have a clear meaning in decision making. Assume that  $\{a_i\}_{i \in I}$  is a sequence of scores assigned to some alternative. The quantity  $\bigotimes_{i \in I} a_i$  is the overall score of the alternative. If the splitting rule is used, the overall score is  $\mathbb{O}$  whenever best and worse scores are opposite. This way of computing the overall score is not very discriminating since many alternatives will get  $\mathbb{O}$  as overall score, even if the scores assigned to them are very different. The two other rules are more discriminating since they discard maximal opposite scores: if best and worst scores are opposite, then look at second best and second worst scores, etc.*

*The purpose of the optimistic and pessimistic rules are merely for illustration of properties. They obviously have no “rational” behaviour in a decision making framework.*

**Remark 7** *The strong rule coincides with the limit of some family of uni-norms proposed by Mesiar and Komorniková [20] (uni-norms are binary operations on  $[0, 1]^2$  which are associative, commutative, non decreasing and with a neutral element  $e \in ]0, 1[$ . See [15] for details).*

Let us endow  $\mathfrak{R}$  with the following order: for  $\langle \cdot \rangle_1, \langle \cdot \rangle_2 \in \mathfrak{R}$ ,  $\langle \cdot \rangle_1 \sqsubseteq \langle \cdot \rangle_2$  iff for all sequences  $\{a_i\}_{i \in I} \in \mathfrak{S}$ ,  $J_1 \supseteq J_2$ , where  $J_1, J_2$  are the index sets of deleted terms for rules 1 and 2.  $\sqsubseteq$  being reflexive, antisymmetric (since computation rules are precisely defined by the set of deleted terms) and transitive,  $(\mathfrak{R}, \sqsubseteq)$  is a partially ordered set. As usual, the interval  $[\langle \cdot \rangle_1, \langle \cdot \rangle_2]$  denotes the set of all computation rules  $\langle \cdot \rangle$  such that  $\langle \cdot \rangle_1 \sqsubseteq \langle \cdot \rangle \sqsubseteq \langle \cdot \rangle_2$ .

Let us introduce another order relation on  $\mathfrak{R}$ . The sequence  $\{a_i\}_{i \in I}$  in  $\mathfrak{S}$  is said to be a *cancelling sequence for the rule  $\langle \cdot \rangle$*  if  $\langle \bigoplus_{i \in I} a_i \rangle = \mathbb{O}$ . We denote by  $\mathcal{O}_{\langle \cdot \rangle}$  the set of cancelling sequences of  $\langle \cdot \rangle$ . We say that computation rule  $\langle \cdot \rangle_1$  is more *discriminating* than rule  $\langle \cdot \rangle_2$ , denoted as  $\langle \cdot \rangle_1 \succ \langle \cdot \rangle_2$ , if  $\mathcal{O}_{\langle \cdot \rangle_1} \subseteq \mathcal{O}_{\langle \cdot \rangle_2}$ . Note that  $\succ$  is only a preorder, since being reflexive and transitive, but not antisymmetric. For a justification of the name “discriminating”, see Remark 3

**Proposition 8** *For any rules  $\langle \cdot \rangle_1, \langle \cdot \rangle_2 \in \mathfrak{R}$ , the following holds:*

- (i)  $\langle \cdot \rangle_1 \sqsubseteq \langle \cdot \rangle_2$  implies that for all sequences  $\{a_i\}_{i \in I} \in \mathfrak{S}$ ,  $|\langle \bigoplus_{i \in I} a_i \rangle_1| \leq |\langle \bigoplus_{i \in I} a_i \rangle_2|$ .
- (ii)  $\langle \cdot \rangle_1 \sqsubseteq \langle \cdot \rangle_2$  implies  $\mathcal{O}_{\langle \cdot \rangle_1} \supseteq \mathcal{O}_{\langle \cdot \rangle_2}$
- (iii)  $\langle \cdot \rangle^+$  is the unique minimal element of  $(\mathfrak{R}, \sqsubseteq)$ , while  $\langle \cdot \rangle_0$  is a maximal element.

**PROOF.** (i) Let  $\{a_i\}_{i \in I}$  not fulfilling associativity. Rule  $k$ ,  $k = 1, 2$ , makes the sequence associative by removing terms  $a_i$ ,  $i \in J_k$ . Then  $\langle \bigoplus_{i \in I \setminus J_k} a_i \rangle_k$  equals either  $\bigvee_{i \in I \setminus J_k} a_i$  or  $\bigwedge_{i \in I \setminus J_k} a_i$ . By hypothesis,  $J_1 \supseteq J_2$ , hence the result.

(ii) Let  $A$  be a cancelling sequence for rule 2, which means that  $\langle \bigoplus_{a \in A} a \rangle_2 = \mathbb{O}$ . Then applying (i), clearly  $A$  is a cancelling sequence for rule 1.

(iii) Obvious for  $\langle \cdot \rangle^+$ .  $\langle \cdot \rangle_0$  is a maximal element since the sequence of deleted terms is  $I_0$ , which is a minimal sequence as remarked above (no proper subset can ensure associativity).  $\square$

The following corollary is immediate.

**Corollary 9**(i) *for any sequence in  $\mathfrak{S}$ ,  $|\langle \bigoplus_{i \in I} a_i \rangle_0| \geq |\langle \bigoplus_{i \in I} a_i \rangle| \geq |\langle \bigoplus_{i \in I} a_i \rangle^+|$ , and  $|\langle \bigoplus_{i \in I} a_i \rangle^+|$  is the lowest bound of  $|\langle \bigoplus_{i \in I} a_i \rangle|$  for all rules  $\langle \cdot \rangle$  in  $\mathfrak{R}$ .*

- (ii)  $\mathcal{O}_{\langle \cdot \rangle_0} \subseteq \mathcal{O}_{\langle \cdot \rangle^+} \subseteq \mathcal{O}_{\langle \cdot \rangle^+}$ .

$(\mathfrak{R}, \sqsubseteq)$  fails to be a lattice or even a semi-lattice, as shown by the following example. Consider the optimistic and pessimistic rules and the following sequence:  $3, 3, 3, 2, 1, -2, -3, -3, -3$ . The terms deleted by these rules are  $J_{\text{opt}} = 3, 3, -3, -3, -3$ , and  $J_{\text{pes}} = 3, 3, 3, -3, -3$ . An upper bound of the two rules deletes at most the terms in  $J_{\text{opt}} \cap J_{\text{pes}} = 3, 3, -3, -3$ . In any case, the resulting sequence is not associative, hence it does not define a computation rule. Similarly, a lower bound deletes at least  $J_{\text{opt}} \cup J_{\text{pes}} = 3, 3, 3, -3, -3, -3$ . It is easy to see that  $3, 3, 3, 2, -3, -3, -3$  and  $3, 3, 3, -2, -3, -3, -3$  are two maximal lower bounds each defining a computation rule, hence there is no greatest lower bound.

We give hereafter some other properties.

**Proposition 10** For any sequences  $\{a_i\}_{i \in I}$ ,  $\{a'_i\}_{i \in I}$ , and  $\{b_i\}_{i \in J}$  in  $\mathfrak{S}$

(i)

$$\bigwedge_{i \in I} a_i \leq \langle \bigotimes_{i \in I} a_i \rangle \leq \bigvee_{i \in I} a_i.$$

(ii) The rules  $\langle \cdot \rangle_{-}^{+}$  and  $\langle \cdot \rangle_0$  are isotone, i.e. they satisfy

$$a_i \leq a'_i, \quad \forall i \in I \text{ implies } \langle \bigotimes_{i \in I} a_i \rangle \leq \langle \bigotimes_{i \in I} a'_i \rangle.$$

(iii)

$$\begin{aligned} |\langle (\bigotimes_{i \in I} a_i) \otimes (\bigotimes_{j \in J} b_j) \rangle_{-}^{+}| &\geq |\langle \bigotimes_{i \in I} a_i \rangle_{-}^{+}| \\ \text{or } \langle (\bigotimes_{i \in I} a_i) \otimes (\bigotimes_{j \in J} b_j) \rangle_{-}^{+} &= \mathbb{O}. \end{aligned}$$

**PROOF.** (i) Clear from definition.

(ii) It suffices to show the result for one argument, say  $a_j$ . Let us consider the rule  $\langle \cdot \rangle_{-}^{+}$ . If  $a_j \geq \mathbb{O}$ , then by Prop. 4 (x),  $\bigotimes_{a_i \geq \mathbb{O}} a_i$  will not decrease when  $a_j$  is replaced by  $a'_j$ , so that  $\bigotimes_{i \in I} a_i$  will not decrease too (similarly if  $a_j < \mathbb{O}$ ).

We turn to the rule  $\langle \cdot \rangle_0$ . We consider the sequence  $\{a_i\}_{i \in I}$ , and the index set of deleted terms  $J$ . If  $j \in I \setminus J$ , then the expression  $\bigotimes_{i \in I \setminus J} a_i$  is isotone provided associativity still holds when  $a_j$  is replaced by  $a'_j$  (see Prop. 4 (x)). Since  $a'_j \geq a_j$ , the only case where associativity is lost is when  $\bigotimes_{i \in I \setminus J} a_i = a_k$  with  $a_k < \mathbb{O}$ , and  $a'_j = -a_k$ . In this case  $a_k, a'_j$  are deleted, and the result is the 2nd largest in absolute value, which is greater or equal to  $a_k$ , hence the rule is still isotone.

Let us consider the case when  $j \in J$ , and suppose that  $\bigotimes_{i \in I \setminus J} a_i = a_k$ . If  $a_j > \mathbb{O}$ , then for  $a'_j > a_j$ , the pair  $(a'_j, -a_j)$  is no more deleted, and the result of computation will be  $a'_j$ . Since  $a'_j > a_j \geq a_k$ , the rule is isotone. Now, if  $a_j < \mathbb{O}$ , for  $a'_j > a_j$ , the pair  $(a'_j, -a_j)$  is no more deleted, and the result becomes  $-a_j$ . Since  $-a_j \geq a_k$ , isotonicity holds in this case too.

(iii) Let us consider a sequence  $\{a_i\}_{i \in I}$ , not fulfilling associativity. We need only to prove the result for a sequence  $\{b_j\}_{j \in J}$  reduced to a singleton  $b_1$ , the general case follows by induction. We denote  $a := \langle \bigotimes_{i \in I} a_i \rangle_{-}^{+}$ , and  $b := \langle (\bigotimes_{i \in I} a_i) \otimes b_1 \rangle_{-}^{+}$ . We have  $a = a^{+} \otimes a^{-}$ , with  $a^{+} := \bigotimes_{a_i \geq \mathbb{O}} a_i$ ,  $a^{-} := \bigotimes_{a_i < \mathbb{O}} a_i$ . Assume that  $a = a^{+}$ . If  $b_1 \geq \mathbb{O}$ , we have  $b = a^{+} \vee b_1 \geq a$ . If  $b_1 < \mathbb{O}$ ,  $b = a$  unless  $b_1 \leq -a^{+}$ . If  $b_1 = -a^{+}$ , then  $b = \mathbb{O}$ , and if  $b_1 < -a^{+}$ , then  $b = b_1$ , so that  $|b| > |a|$ . Assume now that  $a = \mathbb{O}$ , then trivially the result holds. The case where  $a = a^{-}$  works similarly as the case  $a = a^{+}$ .  $\square$

Computation rule  $\langle \cdot \rangle_{-}$  is not isotone, as shown by the following example: take

the sequence  $-3, 3, 1$  in  $\mathbb{Z}$ . Applying the weak rule leads to 1. Now, if 1 is raised to 3, the result becomes 0.

## 5 The ordinal Möbius transform

Throughout this section, let  $(X, \leq)$  denote a locally finite poset (i.e. any segment  $[u, v] := \{x \in X \mid u \leq x \leq v\}$  is finite) with unique minimal element 0. We begin by briefly recalling the classical construction of the Möbius transform (see e.g. [1,23]), and its connection with capacities.

### 5.1 Basic facts on the Möbius transform

Let us consider  $f, g$  two real-valued functions on  $X$  such that

$$g(x) = \sum_{y \leq x} f(y). \quad (15)$$

A fundamental question in combinatorics is to solve this equation, i.e. to recover  $f$  from  $g$ . The solution is given through the *Möbius function*  $\mu(x, y)$  by

$$f(x) = \sum_{y \leq x} \mu(y, x) g(y) \quad (16)$$

where  $\mu$  is defined inductively by

$$\mu(x, y) = \begin{cases} 1, & \text{if } x = y \\ -\sum_{x \leq t < y} \mu(x, t), & \text{if } x < y \\ 0, & \text{otherwise.} \end{cases}$$

More precisely,  $\mu$  is obtained as the inverse of the Riemann function  $\zeta(x, y) := 1$  if  $x \leq y$  and 0 otherwise, in the sense that  $\zeta \star \mu = \delta$ , where  $\star$  is a group operation on real functions on  $X^2$  defined by:

$$(f \star g)(x, y) = \sum_{x \leq u \leq y} f(x, u) g(u, y), \quad x, y \in X,$$

and  $\delta(x, y) = 1$  iff  $x = y$  and 0 otherwise, is the neutral element.

Viewing in equation (15)  $g$  as the primitive function of  $f$ , we can say that in some sense  $f$  is the derivative of  $g$ . Hence,  $\mu(x, y)$  acts as a differential operator.



In the sequel, our main interest will be capacities and set functions, so that the partially ordered set is the Boolean lattice of subsets of a finite set  $N$ , and  $f, g$  are real-valued set functions, or more restrictively capacities. In this case, for any  $A \subseteq B \subseteq N$  we have  $\mu(A, B) = (-1)^{|B \setminus A|}$ , and denoting set functions by  $v, m$ , formulas (15) and (16) become

$$v(A) = \sum_{B \subseteq A} m(B) \quad (17)$$

$$m(A) = \sum_{B \subseteq A} (-1)^{|A \setminus B|} v(B). \quad (18)$$

The set function  $m$  is called the *Möbius transform* of  $v$ . When necessary, we write  $m^v$  to stress the fact it is the Möbius transform of  $v$ . In cooperative game theory,  $m$  is called the *dividend* of the game  $v$  [14,22]. In the field of decision theory,  $v$  is a capacity and its Möbius transform is a fundamental concept (see e.g. Shafer [25], Chateauneuf and Jaffray [3], Grabisch [8]).

## 5.2 The ordinal Möbius transform

Let  $(L, \leq)$  be a linear reflection lattice, with  $L^+$  its positive part. Consider two  $L$ -valued functions on  $X$ , denoted by  $f, g$ , and the formula:

$$g(x) = \langle \bigoplus_{y \leq x} f(y) \rangle. \quad (19)$$

To enforce uniqueness of this expression, we use some rule of computation. By analogy with the classical case, any solution  $f$  to the above equation plays the role of an *ordinal* Möbius transform of  $g$ , defined with respect to the given rule of computation.

Contrary to the classical case, there is not always a solution to this equation, and if there is one, it may be not unique. Consider the following example: take  $X = \{a, b\}$  with  $a < b$ , and  $g(a) = \mathbb{1}$ ,  $g(b) = -\mathbb{1}$ . We necessarily have  $f(a) = g(a) = \mathbb{1}$ , and  $g(b) = f(b) \oplus f(a)$ . But this last equation reads  $-\mathbb{1} = f(b) \oplus \mathbb{1}$ , which is impossible to satisfy. Let us put now  $g(b) = \mathbb{1}$ . Then any  $f$  such that  $f(a) = \mathbb{1}$  and  $f(b) \neq -\mathbb{1}$  is a solution.

The following result shows that, at least for the splitting rule,  $g$  should satisfy some properties.

**Proposition 11** *If Equation (19) has a solution for the splitting rule  $\langle \cdot \rangle_{-}^{+}$ ,*

then necessarily  $g$  fulfills

$$\forall x \in X, \begin{cases} |g(x)| \geq |g(y)|, & \forall y \prec x \\ or \\ g(x) = \mathbb{O}. \end{cases} \quad (*)$$

**PROOF.** Suppose  $(*)$  does not hold. Then there exists some  $x \in X$  such that  $g(x) \neq \mathbb{O}$  and  $|g(x)| < |g(y_0)|$  for some  $y_0 \prec x$ . We have, assuming  $f$  is a solution of (19),

$$\begin{aligned} g(x) &= \langle \bigotimes_{y \leq x} f(y) \rangle_-^+ \\ &= \langle \bigotimes_{y \leq y_0} f(y) \bigotimes \bigotimes_{\substack{y \leq x \\ y \notin [0, y_0]}} f(y) \rangle_-^+ \end{aligned}$$

Applying Prop. 10 (iii), we get:

$$|g(x)| \geq |\langle \bigotimes_{y \leq y_0} f(y) \rangle_-^+| = |g(y_0)| \text{ or } g(x) = \mathbb{O},$$

which contradicts the hypothesis, hence  $f$  is not a solution.  $\square$

In this section, assuming  $|g|$  is isotone (hence fulfilling conditions of Prop. 11), we will give solutions to this equation for a subset of  $\mathfrak{A}$ , which are expressed through the inverse of the Riemann function as in the classical case. Other solutions may exist, but their detailed study is beyond our scope.

We begin by some considerations close to the classical case. We consider the following set of functions:

$$\mathcal{G} = \{f : X^2 \longrightarrow L \mid f(x, x) = \mathbb{1}, \quad f(x, y) = \mathbb{O} \text{ if } x > y\},$$

equipped with the following operation  $\otimes$  internal on  $\mathcal{G}$ :

$$(f \otimes g)(x, y) := \langle \bigotimes_{x \leq u \leq y} [f(x, u) \otimes g(u, y)] \rangle,$$

with the same computation rule as in (19). The  $\otimes$  operation can be defined also when one of the function has domain  $X$ :  $(f \otimes g)(x, y) := \langle \bigotimes_{x \leq u \leq y} [f(u) \otimes g(u, y)] \rangle$ . Contrary to the classical case,  $(\mathcal{G}, \otimes)$  has not the structure of a group. The lack of distributivity in  $(L, \otimes, \otimes)$  forbids the satisfaction of associativity in  $(\mathcal{G}, \otimes)$ . However, a neutral element always exists, and is defined by

$$\delta(x, y) := \begin{cases} \mathbb{1}, & \text{if } x = y \\ \mathbb{O}, & \text{otherwise} \end{cases}$$

as it is easy to check. Left and right inverses of  $f$  may exist and are not unique in general. Specifically, the left inverse  $f^{-1}$  should satisfy:

$$\langle \bigotimes_{x \leq u \leq y} [f^{-1}(x, u) \otimes f(u, y)] \rangle = \begin{cases} \mathbb{1}, & \text{if } x = y \\ \mathbb{0}, & \text{otherwise} \end{cases}$$

from which we deduce that

$$f^{-1}(x, x) = \mathbb{1}, \quad \forall x \in X \quad (20)$$

$$\langle \bigotimes_{x \leq u \leq y} [f^{-1}(x, u) \otimes f(u, y)] \rangle = \mathbb{0}, \quad \forall x < y. \quad (21)$$

Defining  $f^{-1}(x, y) = \mathbb{0}$  whenever  $x > y$  and using (20), we deduce that  $f^{-1}$  belongs to  $\mathcal{G}$ . The following lemma clarifies the situation for the Riemann function  $\zeta(x, y)$ .

**Lemma 12** *The inverse of the Riemann function (left or right) is given by*

$$\begin{aligned} \zeta^{-1}(x, x) &= \mathbb{1}, \quad \forall x \in X \\ \zeta^{-1}(x, y) &= -\mathbb{1}, \quad \forall x, y \in X \text{ such that } x \prec y \end{aligned}$$

for all  $\langle \cdot \rangle \in \mathfrak{R}$ , and for  $x, y$  such that  $x < y$  and  $x \not\prec y$

- For any rule in  $[\langle \cdot \rangle_+^+, \langle \cdot \rangle_+]$ ,  $-\mathbb{1}, \mathbb{0}$  and  $\mathbb{1}$  are possible values for  $\zeta^{-1}(x, y)$ . In particular, if  $\langle \cdot \rangle = \langle \cdot \rangle_+$ , these are the only possible values, and if  $\langle \cdot \rangle = \langle \cdot \rangle_+^+$ , all values in  $L$  are possible.
- There exists no inverse in general for any rule in  $]\langle \cdot \rangle_+, \langle \cdot \rangle_0]$ . If  $X$  is linearly ordered, then  $\zeta^{-1}(x, y) = \mathbb{0}$  is solution for any rule in  $\mathfrak{R}$ .

**PROOF.** We know already from (20) that  $\zeta^{-1}(x, x) = \mathbb{1}$  for any computation rule. Equation (21) for the Riemann function becomes

$$\langle \bigotimes_{x \leq u \leq y} \zeta^{-1}(x, u) \rangle = \mathbb{0}, \quad \forall x < y.$$

If  $x \prec y$ , then clearly we get  $\zeta^{-1}(x, y) = -\mathbb{1}$  as only solution, and for any computation rule. Let us consider  $x, y$  such that  $x \prec u \prec y$ . The above equation reads

$$\langle \mathbb{1} \otimes \bigotimes_{u \prec y} (-\mathbb{1}) \otimes \zeta^{-1}(x, y) \rangle = \mathbb{0}.$$

Note that it suffices to show that the above sequence of terms belongs to  $\mathcal{O}_{\langle \cdot \rangle}$ . In the case of the splitting rule  $\langle \cdot \rangle_+^+$ , clearly any number in  $L$  is solution for  $\zeta^{-1}(x, y)$ . In the case of the weak rule  $\langle \cdot \rangle_+$ , only  $\mathbb{1}, \mathbb{0}, -\mathbb{1}$  are solutions. Then for any rule in  $[\langle \cdot \rangle_+^+, \langle \cdot \rangle_+]$ , the result is proven using Prop. 8 (ii).

Let us consider any rule  $\langle \cdot \rangle$  in  $]\langle \cdot \rangle_-, \langle \cdot \rangle_0]$ . Then there exist some sequences in  $\mathfrak{S}$  for which the index set of deleted terms is strictly included in  $I_-$ . This means that it may exist a poset  $(X, \leq)$  such that the above equation has no solution. Indeed, if  $\zeta^{-1}(x, y) = \mathbb{1}$  or  $-\mathbb{1}$ , the sequence of  $\mathbb{1}, -\mathbb{1}$  we obtain may be such that the index set of deleted terms is strictly included in  $I_-$ , and so the result cannot be  $\mathbb{0}$ . The same holds if  $\zeta^{-1}(x, y)$  takes any other value. In particular, in the case of the strong rule  $\langle \cdot \rangle_0$ , observe that if there is a unique element  $u$  between  $x$  and  $y$ , then  $\zeta_{-1}(x, y) = \mathbb{0}$  is solution (and due to Prop. 8 (ii) (iii), the results extends to any other rule). If there are two elements  $u$  between  $x$  and  $y$ , then  $\zeta_{-1}(x, y) = \mathbb{1}$  is solution. Otherwise, there is no solution.  $\square$

We call *canonical inverse* the solution where  $\zeta^{-1}(x, y) = \mathbb{0}$  when  $x < y$  but  $x \not\prec y$ . It is a solution for all rules in  $[\langle \cdot \rangle_+^-, \langle \cdot \rangle_-]$  (and for any rule in  $\mathfrak{R}$ , if  $X$  is linearly ordered). By extension, we call it *canonical pseudo-inverse* for rules outside  $[\langle \cdot \rangle_+^-, \langle \cdot \rangle_-]$ , when  $X$  is not a linear order. In the sequel we examine under what conditions inverses of the Riemann function permit to build solutions.

If  $(\mathcal{G}, \otimes)$  were a group, then  $g \otimes \zeta^{-1}$  should be solution to the equation. Let us study when  $f = g \otimes \zeta^{-1}$  is indeed a solution. The following is the main result of the paper.

**Theorem 13** *Assume  $g$  is such that  $|g|$  is isotone. Then  $g \otimes \zeta^{-1}$  is solution to Equation (19) for any rule in  $[\langle \cdot \rangle_+^-, \langle \cdot \rangle_-]$ , where  $\zeta^{-1}$  is any inverse of the Riemann function. For rules in  $]\langle \cdot \rangle_-, \langle \cdot \rangle_0]$  and the canonical pseudo-inverse,  $g \otimes \zeta^{-1}$  is not a solution in general.*

(see proof in Appendix)

Equation (19) may have no solution at all for the strong rule, even if  $|g|$  is isotone. Indeed, take  $X = \{0, a, b, c\}$  with  $0 \prec a \prec c$  and  $0 \prec b \prec c$ ,  $b, c$  being incomparable, and define  $g(0) = \mathbb{0}$ ,  $g(a) = g(b) = -\mathbb{1}$ , and  $g(c) = \mathbb{1}$ . Then clearly  $f(0) = \mathbb{0}$ ,  $f(b) = f(c) = -\mathbb{1}$  and there is no solution for  $f(c)$ .

The preceding results can be summarized as follows.

**Summary 1** *We consider  $f, g : X \longrightarrow L$ , and the following equation to solve:*

$$g(x) = \left\langle \bigotimes_{y \leq x} f(y) \right\rangle$$

*with  $\langle \cdot \rangle \in [\langle \cdot \rangle_+^-, \langle \cdot \rangle_-]$ . We call Möbius function  $\mu(x, y)$  any inverse  $\zeta^{-1}$  of the Riemann function, as given in Lemma 12, and call canonical Möbius function the canonical inverse of the Riemann function.*

*Assuming that  $|g|$  is isotone, then  $f(x) = (g \otimes \mu)(x)$  is solution for any Möbius*

function, where  $\circledast$  is defined with respect to the corresponding computation rule. We call any such  $f$  a Möbius transform of  $g$ , and canonical Möbius transform of  $g$ , denoted  $m^g$ , the one corresponding to the canonical Möbius function. It is given by:

$$m^g(x) := \langle g(x) \circledast \left[ - \bigvee_{y \prec x} g(y) \right] \rangle. \quad (22)$$

### 5.3 The case of non negative isotone functions

A particular case of interest is to restrict to isotone functions valued on  $L^+$  (capacities correspond to this case, hence its interest). Let us call them non negative isotone functions.

**Theorem 14** *For any non negative isotone function  $g$ , the set of non negative solutions to Equation (19) is the interval  $[m_*, m^*]$ , defined by:*

$$\begin{aligned} m^*(x) &= g(x), \quad \forall x \in X \\ m_*(x) &= m^g(x) = \begin{cases} g(x), & \text{if } g(x) > g(y), \quad \forall y \prec x \\ \mathbb{O}, & \text{otherwise} \end{cases}, \quad \forall x \in X. \end{aligned}$$

**PROOF.** Since  $g$  is isotone and non negative,  $m^*$  is clearly a solution. On the other hand, Th. 13 applies, and we recognize  $m_*$  as the canonical Möbius transform (22).

We have to prove that these are indeed the lower and upper bounds of non negative solutions. If  $m^*$  were not the upper bound, it should exist  $x_0 \in X$  such that  $m^*(x_0) > g(x_0)$ . Then due to isotonicity, we would have  $g(x_0) < \bigvee_{y \leq x_0} m^*(y)$ , a contradiction. Similarly, if  $m_*(x_0) < g(x_0)$  for some  $x_0$  such that  $g(x_0) > g(y) > \mathbb{O}$  for all  $y \prec x_0$ , we would have  $g(x_0) > \bigvee_{y \leq x} m_*(y)$ , a contradiction again.

Lastly, we show that any  $f \in [m_*, m^*]$  is also a solution. Since  $m_*, m^*$  are non negative solutions, we have for any  $x$

$$\bigvee_{y \leq x} m_*(y) = \bigvee_{y \leq x} m^*(y).$$

Since  $\bigvee$  is increasing, any  $m \in [m_*, m^*]$  will also satisfy the equation.  $\square$

In case of no ambiguity, we denote simply  $m^g$  by  $m$ . Moreover, since our framework is ordinal in the rest of the paper, we will omit to call it “ordinal”,

and will use the term “classical” Möbius transform when referring to the usual definition. We denote by  $[m]$  the interval  $[m_*, m^*]$ , and with some abuse of notation, any function in this interval.

Some remarks are of interest at this stage.

**Remark 15** *As with the classical case, the Möbius transform has the meaning of a derivative. From Definition 1 and (11), it is clear that  $m \equiv g'$ .*

**Remark 16** *Since  $f, g$  are non negative, we need no more computation rules in (19). However, negative solutions exist. It is easy to check that for any computation rule,  $m_*$  can be defined by*

$$m_*(x) = \begin{cases} g(x), & \text{if } g(x) > g(y), \quad \forall y \prec x \\ \text{any } e \in L, e \succ -g(x), & \text{otherwise} \end{cases},$$

$\forall x \in X$ . However, negative solutions do not possess good properties, and would not permit to obtain the subsequent results.

**Definition 17** *Let  $g$  be any isotone function from  $X$  to  $L^+$ . We call  $g$ -chain any chain  $C$  in  $X$  such that  $g(x)$  is constant on  $C$ , and there is no chain  $C' \supset C$  such that  $g(x)$  is constant on  $C'$ . The set of all  $g$ -chains is denoted  $\mathcal{C}(g)$ . The value of a  $g$ -chain  $C$  is defined by  $g(C) := g(x)$  for some  $x \in C$ .*

Any  $g$ -chain  $C$  has a unique minimal element, denoted  $C_*$ . Indeed, either  $C$  is finite or infinite. In the first case, the results trivially hold. In the second case, since 0 is the unique minimal element of  $X$ , and  $X$  is locally finite,  $C$  has the form  $\{x | x \geq a\}$ , hence the result. On the contrary, there is not always a maximal element  $C^*$ .

If a  $g$ -chain  $C$  is finite, its length is defined as usual by  $l(C) := |C| - 1$ .

The following is easy to show (proof is omitted).

**Proposition 18** *Let  $g$  be any isotone function from  $X$  to  $L^+$ , and  $C$  be any  $g$ -chain. Then:*

- (i)  $\mathcal{C}(g) = \emptyset$  iff  $m \equiv g$ .
- (ii) If  $(X, \leq) = (2^N, \subseteq)$  where  $N$  is a finite set of  $n$  elements, and  $g(\emptyset) < g(N)$ , then  $l(C) < n$  (i.e.  $C$  is not a maximal chain), for any  $C \in \mathcal{C}(g)$ .
- (iii) Let  $\mathcal{C}(g) \neq \emptyset$  and  $C \in \mathcal{C}(g)$ . Then
  - (iii.1) For all  $x \in C$ ,  $x \neq C_*$ ,  $m(x) = \mathbb{O}$ .
  - (iii.2)  $m(C_*) = g(C_*)$ .
  - (iii.3) For all  $x \notin C$ ,  $\forall C \in \mathcal{C}(g)$ ,  $m(x) = g(x)$ .

Let us suppose now that  $X$  and  $L^+$  are endowed with a conjugation mapping  $\bar{\cdot}$ . Then necessarily,  $X$  has a unique maximal element, denoted 1, and any  $g$ -chain is finite, with a unique maximal and minimal element. We define the *conjugate of  $g$*  by  $\bar{g}(x) := \overline{g(\bar{x})}$ . In this section, we compute  $m^{\bar{g}}$  and express it with respect to  $m^g$ . The following can be shown.

**Proposition 19** *Under the above assumptions, let  $\bar{g}$  be the conjugate function of  $g : X \longrightarrow L^+$ . Then:*

(i) *the set of  $\bar{g}$ -chains is given by*

$$\mathcal{C}(\bar{g}) = \{\bar{C} := \{\bar{c}_l, \dots, \bar{c}_1\} \mid \{c_1, \dots, c_l\} =: C, C \in \mathcal{C}(g)\}.$$

$$\text{and } \bar{g}(\bar{C}) = \overline{g(C)}.$$

(ii) *the Möbius transform of  $\bar{g}$  is given by*

$$m^{\bar{g}}(x) = \begin{cases} \mathbb{O}, & \text{for all } x \text{ in some } \bar{C} \in \mathcal{C}(\bar{g}), x \neq \bar{C}_* \\ \overline{m^g(C_*)}, & \text{if } x = \bar{C}_* \\ \overline{m^g(\bar{x})}, & \text{otherwise.} \end{cases} \quad (23)$$

**PROOF.** (i) Let us consider  $C \in \mathcal{C}(g)$ , and  $C = \{c_1, \dots, c_l\}$  with  $c_1 < \dots < c_l$ . Since  $g$  is constant over  $C$ , we get  $\bar{g}(\bar{c}_1) = \dots = \bar{g}(\bar{c}_l)$ ,  $\bar{c}_l < \dots < \bar{c}_1$ , which means that  $\bar{C} := \{\bar{c}_l, \dots, \bar{c}_1\}$  is a  $\bar{g}$ -chain. Also, we have  $g(C) = \overline{g(\bar{C})}$ .

(ii) Suppose that  $\mathcal{C}(g) = \emptyset$ . Then  $\mathcal{C}(\bar{g}) = \emptyset$  too, and due to Prop. 18 (i),  $m^{\bar{g}} \equiv \bar{g}$ . This leads to

$$m^{\bar{g}}(x) = \overline{m^g(\bar{x})}.$$

Suppose now that  $\mathcal{C}(g) \neq \emptyset$ , and  $C \in \mathcal{C}(g)$ , with corresponding  $\bar{C} \in \mathcal{C}(\bar{g})$ . By Prop. 18 (iii), if  $c \in \bar{C}$ ,  $c \neq \bar{C}_*$ , then  $m^{\bar{g}}(c) = \mathbb{O}$ . If  $c = \bar{C}_*$ , then  $m^{\bar{g}}(c) = \bar{g}(c) = \overline{g(\bar{c})}$ . Remark that  $\bar{c} = C^*$ , so that  $m^g(\bar{c}) = \mathbb{O}$ . But  $m^g(C_*) = g(C_*) = g(\bar{c})$  (since  $\bar{c}$  and  $C_*$  are in  $C$ ), hence the result.  $\square$

It is possible to have a slightly more compact form for this result. Let us denote by  $n_C(c)$  the element in  $C$  which has the symmetric place of  $c$  (i.e.  $n_C(c_k) = c_{l-k+1}$ ). We can write for  $c = \bar{C}_*$ :

$$m^{\bar{g}}(c) = \overline{m^g(n_C(\bar{c}))}$$

(see figure 3 below). Considering that any  $c$  not belonging to a  $g$ -chain is itself a chain  $C$  of 0 length, so that  $n_C(c) = c$ , we have the general result:

$$m^{\bar{g}}(x) = \begin{cases} \mathbb{O}, & \text{for all } x \text{ in some } \bar{C} \in \mathcal{C}(\bar{g}), x \neq \bar{C}_* \\ \overline{m^g(n_C(\bar{x}))}, & \text{otherwise.} \end{cases} \quad (24)$$

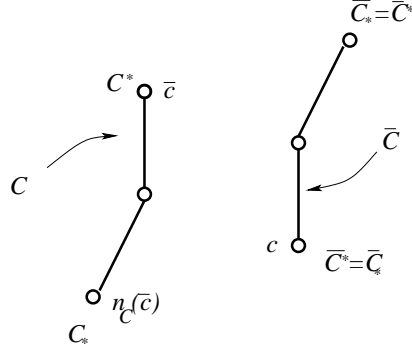


Fig. 3.  $g$ -chains and  $\bar{g}$ -chains

The classical Möbius transform can be viewed as a linear operator on the set of real functions on  $X$ . We may expect that the ordinal counterpart has a similar property with  $\odot, \oslash$ , i.e.  $m^{f \odot g} = m^f \oslash m^g$  and  $m^{\alpha \odot f} = \alpha \oslash m^f$ , for any  $\alpha \in L^+$ . However, the following simple example shows that this is not the case.

**Example 20** *Let us take  $X$  to be the Boolean lattice  $2^2$  whose elements are denoted  $\emptyset, \{1\}, \{2\}, \{1, 2\}$ , and consider two functions  $g_1, g_2$  defined as follows:*

	$\emptyset$	$\{1\}$	$\{2\}$	$\{1, 2\}$
$g_1$	$\oslash$	$\oslash$	$\oslash$	$\mathbb{1}$
$g_2$	$\oslash$	$\mathbb{1}$	$\mathbb{1}$	$\mathbb{1}$

*The computation of the Möbius transform  $m_*$  gives*

	$\emptyset$	$\{1\}$	$\{2\}$	$\{1, 2\}$
$m_*[g_1]$	$\oslash$	$\oslash$	$\oslash$	$\mathbb{1}$
$m_*[g_2]$	$\oslash$	$\mathbb{1}$	$\mathbb{1}$	$\oslash$

*Clearly,  $g_1 \odot g_2 = g_2$ , but  $m_*^{g_1} \oslash m_*^{g_2} \neq m_*^{g_2}$ .*

Remarking that  $m^*$  is maxitive, one should expect that it is possible to find some  $m \in [m], m < m^*$  at least on some element of  $X$ . The above example shows that this is even impossible in general: due to the fact that  $m^{g_1}(\{1, 2\}) = \mathbb{1}$ , we must have  $m^{g_2}(\{1, 2\}) = \mathbb{1}$ , and thus  $m \equiv m^*$ .



## 6 The ordinal Möbius transform of capacities

We devote this section to the particular case of capacities on some finite set  $N := \{1, \dots, n\}$ , which is our original motivation in this paper. Then  $X = 2^N$  is a Boolean lattice, and we suppose in addition that  $L^+$  is a conjugation linear lattice. Capacities are denoted by  $v$ .

A first fact is that we can give an alternative expression of the (canonical) Möbius transform, which is very similar to the classical one (18).

$$m(A) := \bigvee_{B \subseteq A, |A \setminus B| \text{ even}} v(B) \oslash \left( - \bigvee_{B \subseteq A, |A \setminus B| \text{ odd}} v(B) \right) \quad (25)$$

for any  $A \subseteq N$ . Indeed,

$$\bigvee_{B \subseteq A, |A \setminus B| \text{ even}} v(B) = v(A)$$

and

$$\bigvee_{B \subseteq A, |A \setminus B| \text{ odd}} v(B) = \bigvee_{B \prec A} v(B)$$

so that we recognize (22).

In the field of decision theory and artificial intelligence, sup-preserving functions from  $X$  to  $L^+$  (i.e. such that  $g(x \vee y) = g(x) \vee g(y)$ , for every  $x, y \in X$ ) are called *possibility measures* [29,6] or *maxitive measures*, and are denoted by  $\Pi$ . By conjugation we have  $\bar{g}(x \wedge y) = \bar{g}(x) \wedge \bar{g}(y)$ , for every  $x, y \in X$  (inf-preserving functions), they are called *necessity measures* or *minitive measures*, and are denoted by  $N$ . Remark that for any  $A = \{i_1, \dots, i_l\} \subseteq N$ , we have  $\Pi(A) = \bigvee_{i \in A} \Pi(\{i\})$ . The following can be shown.

**Theorem 21** *Let  $\Pi, N$  be a pair of conjugate possibility and necessity measures, and suppose without loss of generality that the elements in  $N$  are such that  $\Pi(\{1\}) \leq \dots \leq \Pi(\{n\})$ . Then*

- *the Möbius transform of  $\Pi$  is non zero on an antichain:*

$$m^\Pi(A) = \begin{cases} \Pi(\{i\}), & \text{if } A = \{i\}, i \in N \\ \mathbb{O}, & \text{otherwise.} \end{cases}$$

- *the Möbius transform of  $N$  is non zero on a chain. Assuming  $\mathbb{O} < \Pi(\{1\}) < \dots < \Pi(\{n\}) = \mathbb{1}$ , the expression is:*

$$m^N(A) = \begin{cases} \overline{\Pi(\{i\})}, & \text{if } A = \{i+1, \dots, n\}, i \in N \\ \mathbb{O}, & \text{otherwise.} \end{cases}$$

If  $\Pi(\{i\}) = \Pi(\{i+1\})$  for some  $i$ , then  $m^N(\{i+1, \dots, n\}) = \mathbb{O}$ .

**PROOF.** Let us suppose  $\mathbb{O} < \Pi(\{1\}) < \dots < \Pi(\{n\}) = \mathbb{1}$ .

(i) Let us compute  $\mathcal{C}(\Pi)$ . Let us consider  $i \in N$ . We denote by  $L_i$  the sublattice which is the interval  $[\{i\}, \{1, \dots, i\}]$ . By construction, any subset  $A \in L_i$  is such that  $\Pi(A) = \Pi(\{i\})$ , and only those ones, which proves that all  $\Pi$ -chains with value  $\Pi(\{i\})$  are the maximal chains of  $L_i$ . In other words,  $\mathcal{G}(\Pi(\{i\})) = L_i$ . Now, the bottom of  $L_i$  being  $\{i\}$ , we get the result.

(ii) From Prop. 19 (i), we know that  $\mathcal{C}(N)$  is in some sense the symmetric of  $\mathcal{C}(\Pi)$  in the lattice  $2^N$ . More precisely, the sublattices of interest are  $\bar{L}_i := [\{1, \dots, i\}, \{i\}]$ . They correspond to the groups  $\mathcal{G}(\overline{\Pi}(\{i\}))$ , and since the bottom element of  $\bar{L}_i$  is  $\{1, \dots, i\}$ , and only  $N$  does not belong to any  $\bar{L}_i$ , we get the desired result.

If  $\Pi(\{i\}) = \Pi(\{i+1\})$  for some  $i$ , then it is easy to check that subset  $\{i+1, \dots, n\}$  disappears in the chain, but there is no change for  $\Pi$ .  $\square$

Figure 4 illustrates the result.

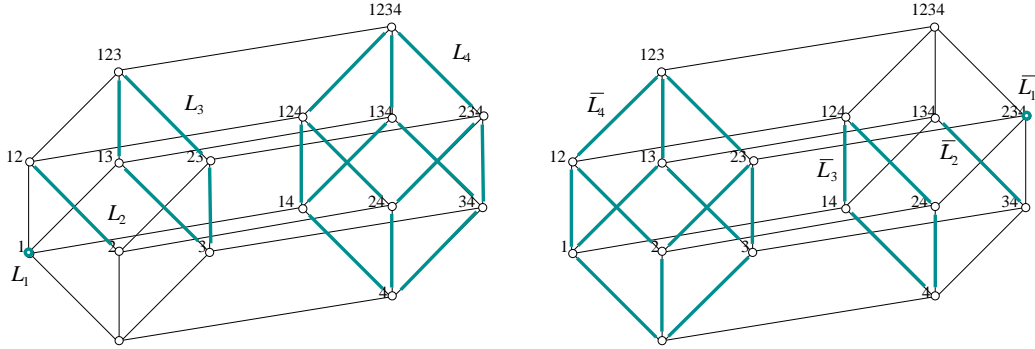


Fig. 4.  $\Pi$ -chains (left) and  $N$ -chains (right) with  $N = \{1, 2, 3, 4\}$ .

## 7 Applications of symmetric ordered structures and perspectives

We conclude the paper by indicating several possible applications of our symmetric ordered structure. We mainly developed in this paper the theory of Möbius transform, and its application to capacities. We briefly mentioned in the introduction that one of the main motivation was the definition of a symmetric Sugeno integral. Clearly, our aim is achieved, since Equation (3) is now perfectly defined, and could be a starting point to develop an ordinal or

*qualitative* counterpart of Cumulative Prospect Theory, a theory which is of primary importance in e.g. economics. We refer the reader to [10] for a detailed study of the symmetric Sugeno integral, along with other results on capacities and the ordinal Möbius transform. Based on symmetric ordered structures and the symmetric Sugeno integral, we have already built a general model of multicriteria decision making [11], which permits to tackle real problems where only qualitative information is available. This is indeed a common situation in many applications (e.g. project selection, subjective evaluation of consumer goods, etc.).

Another application would be to investigate capacities defined on arbitrary lattices instead of the usual Boolean lattice [12], a new promising topic in decision making. Considering these general capacities, valued on  $L$  instead of a real interval, we need our general results from Sections 4 and 5 to get the Möbius transform and properly define a general Sugeno integral.

On a purely mathematical point of view, we have studied in detail algebraic properties of our new structure, and in particular the possible ways to escape from non associativity. The generality of our results may open new areas related to ordered structures and combinatorics. It might also be viewed as a starting point of ordinal “linear” algebra, noticing that  $\otimes$  is in fact the matrix product. We describe hereafter a possible application of this ordinal linear algebra. Considering two finite universal sets  $X, Y$ , a *fuzzy binary relation* or valued binary relation on  $X \times Y$  is simply a function  $R : X \times Y \longrightarrow [0, 1]$ , where  $R(x, y)$  is the strength of relation between  $x$  and  $y$ . Many results exist in this area (see e.g. [21, 7]), but we are interested here in what is called *fuzzy relation equations*, which are important in system theory. Considering finite universal sets  $X, Y, Z$  and three fuzzy relations  $P, Q, R$  on  $X \times Y, Y \times Z, X \times Z$  respectively, we consider the equation  $R = P \circ Q$ , which we want to solve for  $P$ . Composition of relations is given by:

$$R(x, z) = \bigvee_{y \in Y} (P(x, y) \wedge R(y, z)).$$

The solution set of this equation, whenever non empty, has the structure of a union of intervals  $[\check{P}_i, \hat{P}]$ , where  $\hat{P}$  is the unique maximal solution, and  $\check{P}_i$  are minimal ones. Allowing fuzzy relations to be valued in  $[-1, 1]$  or any symmetric linear order (*bipolar* fuzzy relation), replacing  $\vee, \wedge$  by  $\oplus, \ominus$  and considering a particular computation rule, the above equation coincides with our  $\otimes$  operation. Hence our results could provide powerful tools for solving bipolar fuzzy relation equations, a topic which has never been addressed, but which may become important in the near future, since bipolar scales deserve a great interest in this field.

An interesting further study would be to change the starting point, e.g. en-

forcing associativity and loosing symmetry. We already know from Prop. 5 (2) some properties of this kind of structure.

## Acknowledgements

This work has benefited from many fruitful discussions with D. Denneberg, J.L. Marichal, and T. Murofushi, who are deeply acknowledged. The author thanks also the anonymous referees for their constructive comments.

## References

- [1] C. Berge. *Principles of Combinatorics*. Academic Press, 1971.
- [2] G. Birkhoff. *Lattice Theory*. American Mathematical Society, 3d edition, 1967.
- [3] A. Chateauneuf and J.Y. Jaffray. Some characterizations of lower probabilities and other monotone capacities through the use of Möbius inversion. *Mathematical Social Sciences*, 17:263–283, 1989.
- [4] G. Choquet. Theory of capacities. *Annales de l'Institut Fourier*, 5:131–295, 1953.
- [5] D. Denneberg and M. Grabisch. Functions and measures with linear ordinal scales, (a-)symmetric Sugeno integral and ordinal Ky Fan distance. In *2nd Int. Symposium on Imprecise Probabilities and Their Applications*, Ithaca, New York, June 2001.
- [6] D. Dubois and H. Prade. *Fuzzy Sets and Systems: Theory and Applications*. Academic Press, 1980.
- [7] J.C. Fodor and M. Roubens. *Fuzzy Preference Modelling and Multi-Criteria Decision Aid*. Kluwer Academic Publisher, 1994.
- [8] M. Grabisch.  $k$ -order additive discrete fuzzy measures and their representation. *Fuzzy Sets and Systems*, 92:167–189, 1997.
- [9] M. Grabisch. On the representation of  $k$ -decomposable measures. In *7th IFSA World Congress*, pages 478–483, Prague, Czech Republic, June 1997.
- [10] M. Grabisch. The symmetric Sugeno integral. *Fuzzy Sets and Systems*, 139:473–490, 2003.
- [11] M. Grabisch, S. Dia, and Ch. Labreuche. A multicriteria decision making framework in ordinal context based on Sugeno integral. In *Joint 9th IFSA World Congress and 20th NAFIPS Int. Conf.*, Vancouver, Canada, July 2001.

- [12] M. Grabisch and Ch. Labreuche. Interaction between attributes in a general setting for knowledge discovery. In *4th Int. JIM Conf. (Journées de l'Informatique Messine) on Knowledge Discovery and Discrete Mathematics*, pages 215–222, Metz, France, September 2003.
- [13] G. Grätzer. *General Lattice Theory*. Birkhäuser, 2nd edition, 1998.
- [14] J.C. Harsanyi. A simplified bargaining model for the  $n$ -person cooperative game. *International Economic Review*, 4:194–220, 1963.
- [15] E.P. Klement, R. Mesiar, and E. Pap. *Triangular Norms*. Kluwer Academic Publishers, Dordrecht, 2000.
- [16] J.L. Marichal. On Sugeno integral as an aggregation function. *Fuzzy Sets and Systems*, 114:347–365, 2000.
- [17] J.L. Marichal. Personal communication, 2001.
- [18] J.L. Marichal, P. Mathonet, and E. Tousset. Mesures floues définies sur une échelle ordinale. working paper, 1996.
- [19] R. Mesiar.  $k$ -order pan-additive discrete fuzzy measures. In *7th IFSA World Congress*, pages 488–490, Prague, Czech Republic, June 1997.
- [20] R. Mesiar and M. Komorníková. Triangular norm-based aggregation of evidence under fuzziness. In B. Bouchon-Meunier, editor, *Aggregation and Fusion of Imperfect Information*, Studies in Fuzziness and Soft Computing, pages 11–35. Physica Verlag, 1998.
- [21] A. Di Nola, S. Sessa, W. Pedrycz, and E. Sanchez. *Fuzzy relational equations and their applications to knowledge engineering*. Kluwer Academic, 1989.
- [22] G. Owen. Multilinear extensions of games. In A.E. Roth, editor, *The Shapley Value. Essays in Honor of Lloyd S. Shapley*, pages 139–151. Cambridge University Press, 1988.
- [23] G.C. Rota. On the foundations of combinatorial theory I. Theory of Möbius functions. *Zeitschrift für Wahrscheinlichkeitstheorie und Verwandte Gebiete*, 2:340–368, 1964.
- [24] D. Schmeidler. Integral representation without additivity. *Proc. of the Amer. Math. Soc.*, 97(2):255–261, 1986.
- [25] G. Shafer. *A Mathematical Theory of Evidence*. Princeton Univ. Press, 1976.
- [26] M. Sugeno. *Theory of fuzzy integrals and its applications*. PhD thesis, Tokyo Institute of Technology, 1974.
- [27] A. Tversky and D. Kahneman. Advances in prospect theory: cumulative representation of uncertainty. *J. of Risk and Uncertainty*, 5:297–323, 1992.
- [28] S. Weber.  $\perp$ -decomposable measures and integrals for archimedean t-conorms  $\perp$ . *J. Math. Anal. Appl.*, 101:114–138, 1984.
- [29] L.A. Zadeh. Fuzzy sets as a basis for a theory of possibility. *Fuzzy Sets & Systems*, 1:3–28, 1978.

## A Proof of Theorem 13

We first need a technical lemma.

**Lemma 22** *Let  $f = g \otimes \zeta^{-1}$ , where  $\zeta^{-1}$  is any inverse of the Riemann function, and assume that  $|g|$  is isotone. Then, for any  $x \in X$ :*

(i) *If for all  $y \prec x$ ,  $|g(x)| > |g(y)|$  or  $g(x) = -g(y)$ ,*

$$f(x) = g(x)$$

*for any computation rule in  $\mathfrak{R}$ .*

(ii) *If there exists  $y \prec x$  such that  $g(x) = g(y)$*

(ii.1) *in the case of the splitting rule  $\langle \cdot \rangle_+^+$ , then  $f(x) = \mathbb{O}$ .*

(ii.2) *for any rule in  $[\langle \cdot \rangle_-^+, \langle \cdot \rangle_-]$ , we have  $|f(x)| < |g(x)|$ .*

(iii) *For any rule in  $[\langle \cdot \rangle_-, \langle \cdot \rangle_0]$ , and for the canonical pseudo-inverse*

$$f(x) = \begin{cases} \mathbb{O}, & \text{if } |G^+| = |G^-| + 1 \\ \text{either } g(x), \mathbb{O} \text{ or } -g(x) & \text{otherwise} \end{cases}$$

*with  $G^+ := \{y \in X; y \prec x \text{ and } g(y) = g(x)\}$ , and  $G^- := \{y \in X; y \prec x \text{ and } g(y) = -g(x)\}$ . In the case of the strong rule  $\langle \cdot \rangle_0$ , the result particularizes as follows*

$$f(x) = \begin{cases} \mathbb{O}, & \text{if } |G^+| = |G^-| + 1 \\ g(x), & \text{if } |G^+| \leq |G^-| \\ -g(x), & \text{otherwise.} \end{cases}$$

(iv) *if  $y \leq x$  and  $|g(y)| < |g(x)|$ , then  $|f(y)| < |g(x)|$  for all computation rules in  $\mathfrak{R}$ .*

**PROOF.** We have

$$\begin{aligned} f(x) &= g \otimes \zeta^{-1}(x) = \langle \bigotimes_{u \leq x} g(u) \otimes \zeta^{-1}(u, x) \rangle \\ &= \langle g(x) \otimes \bigotimes_{u \prec x} (-g(u)) \otimes \bigotimes_{u \prec v \prec x} (g(u) \otimes \zeta^{-1}(u, x)) \otimes \dots \rangle. \end{aligned} \quad (\text{A.1})$$

Let us remark that

$$|g(u) \otimes \zeta^{-1}(u, x)| \leq |g(u)| \quad (\text{A.2})$$

for all  $u \leq x$ . If  $|g(x)| > |g(y)|$  for all  $y \prec x$  or  $g(x) = -g(y)$  for some  $y \prec x$ , then by (A.2) clearly associativity holds in (A.1), so that for any computation rule the result is the same, which is  $g(x)$ . This proves (i).

Suppose there is some  $y \prec x$  such that  $g(x) = g(y)$ . Using (A.2) and due to the isotonicity of  $|g|$ , extremal terms in  $f(x)$  are  $g(x)$  and  $-g(x)$ . This proves

that  $f(x) = \mathbb{O}$  for the splitting rule. If the weak rule is used, then all terms  $g(x), -g(x)$  disappear, so that we can only deduce that  $|f(x)| < |g(x)|$ . Now, using Prop. 8 (i), we have proven (ii).

(iii) is clear since  $f(x) = \langle g(x) \otimes (-\otimes_{y \prec x} g(y)) \rangle$ . (iv) comes from (A.2), isotonicity of  $|g|$ , and Prop. 10 (i).  $\square$

**PROOF.** (Th. 13) We assume that  $g(x) \neq \mathbb{O}$ , otherwise the result holds trivially. Let  $x \in X$ . Assume  $|g(x)| > |g(y)|$  or  $g(x) = -g(y)$  for all  $y \prec x$ . Then by Lemma 22 (i),  $f(x) = g(x)$ , and by Lemma 22 (iv),  $|f(y)| < |g(x)|$ . Hence  $\langle \otimes_{y \leq x} f(y) \rangle = \langle f(x) \otimes \otimes_{y < x} f(y) \rangle = g(x)$  as expected, since associativity holds.

Assume  $g(x) = g(y)$  for some  $y \prec x$ . Let us introduce  $C_x := \{y \in X | g(y) = g(x), y < x\}$ . Since 0 is the unique minimal element of  $X$ ,  $C_x \subseteq [0, x]$  and hence is finite. Thus,  $C_x$  possesses at least one minimal element. Let us denote by  $C_{x*}$  the set of these minimal elements. We have:

$$\langle \otimes_{y \in C_x} f(y) \rangle = \langle \otimes_{y \in C_{x*}} f(y) \otimes \otimes_{y \in C_x \setminus C_{x*}} f(y) \rangle.$$

From Lemma 22 (i), we have  $f(y) = g(y) = g(x)$  for all  $y \in C_{x*}$ , and for all  $y \in C_x \setminus C_{x*}$ , we have  $|f(y)| < |g(y)|$  for any rule in  $[\langle \cdot \rangle_+^+, \langle \cdot \rangle_-]$  (use Lemma 22 (ii) and the fact that  $|g(x)| > \mathbb{O}$ ). Hence  $\langle \otimes_{y \in C_x} f(y) \rangle = g(x)$  since associativity holds. Now,

$$\langle \otimes_{y \leq x} f(y) \rangle = \langle \otimes_{y \in C_x} f(y) \otimes \otimes_{\substack{y < x \\ y \notin C_x}} f(y) \rangle.$$

Since  $\otimes_{y \in C_x} f(y) = g(x)$  and by Lemma 22 (iv)  $|f(y)| < |g(x)|$  for all  $y < x, y \notin C_x$ , we have finally  $\langle \otimes_{y \leq x} f(y) \rangle = g(x)$  as desired.

Consider now a rule in  $[\langle \cdot \rangle_-, \langle \cdot \rangle_0]$  and the canonical pseudo-inverse. From Lemma 22 (iii), we may have  $f(y) = -g(x)$  for some  $y \in C_x \setminus C_{x*}$ , so that  $\langle \otimes_{y \in C_x} f(y) \rangle \neq g(x)$  may occur, and the result does not hold.  $\square$